Statistical Approach to NoC Design

Itamar Cohen, Ori Rottenstreich and Isaac Keslassy

Technion (Israel)



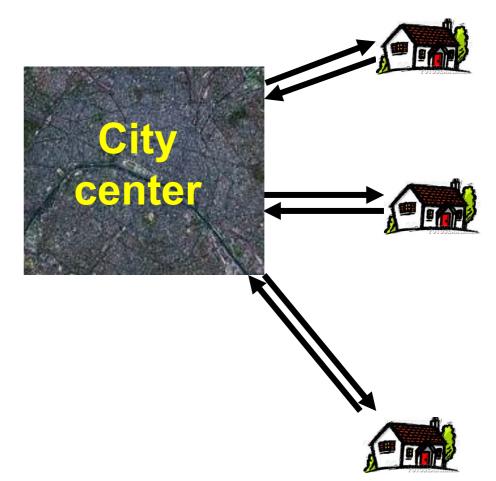


The traffic matrix in NoCs is often-changing and unpredictable

⇒ makes NoCs hard to design

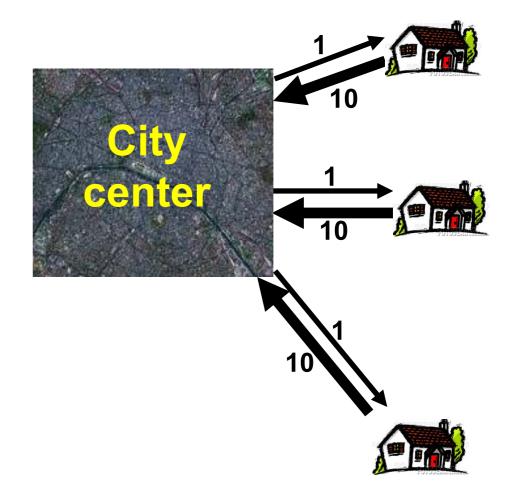
Road Capacities

- Goal: Design road capacities between a city and its suburbs
- Let's model the traffic matrices...



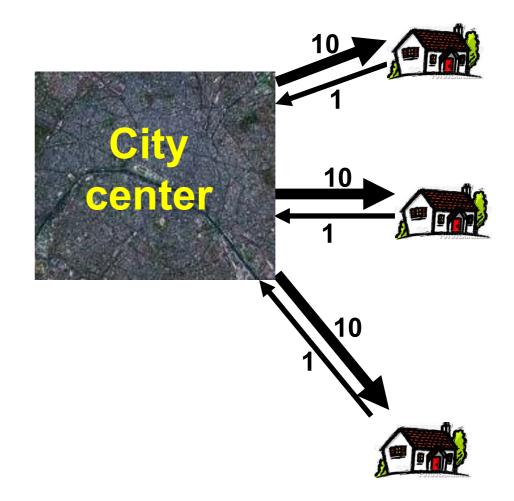
Road Capacities

Morning peak: most traffic towards city



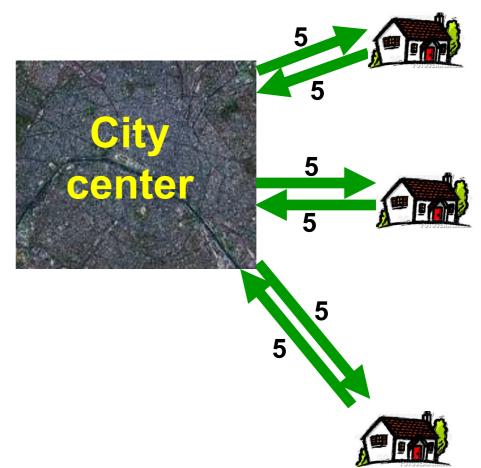
Road Capacities

- Morning peak: most traffic towards city
- Afternoon peak: most traffic leaving city



Solution (1): Average-Case

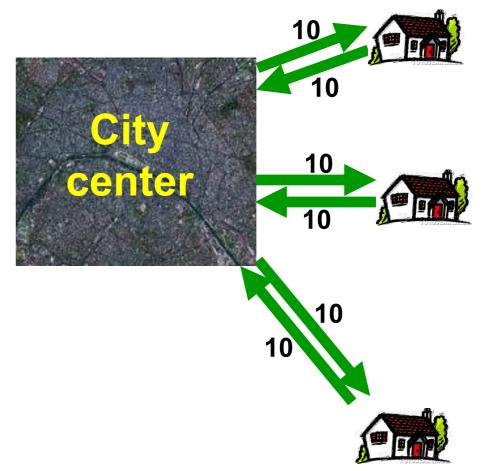
- Solution (1): plan for average-case
 - i.e. allocate capacity of ~5 for each link.
 - > λ < μ
- Problem: traffic jam during many hours, every day.



Solution (2): Worst-Case

Solution (2): plan for worst-case

 i.e. allocate capacity of 10 for each link.

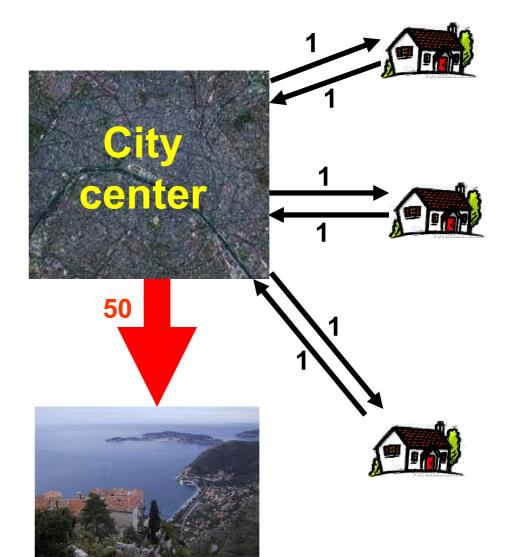


Holidays

 Problem: traffic burst in the holidays
 ⇒ excessive resources

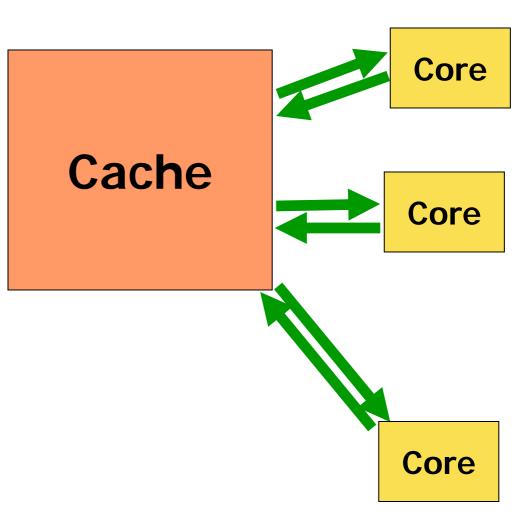
Solution (3): statistical approach

- Enough capacity for 99% of the time
- Allow for occasional congestion



Back to the NoC world

- Similar problems in NoC design process
 - > City \rightarrow Shared cache
 - > Suburbs \rightarrow Cores
 - Many possible traffic matrices: writing, reading, etc.



Motivation

Tradeoff: more resources vs. better guarantee

"Easy" solution: Worst-Case Guarantee

E.g., no-congestion guarant time

Better when no congestion is allowed

Objective: Statistical Guarantee

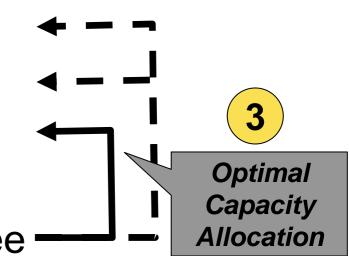
E.g., no-congestion guaranteet time

Better when power/area are most important

Statistical Approach to NoC Design

Given:

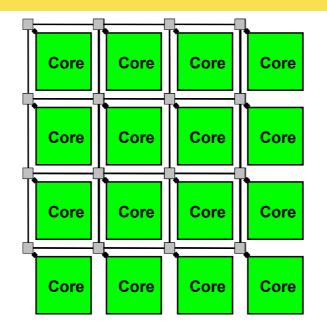
- Traffic matrix distribution
 - Topology
 - Routing
 - Link capacities



Compute congestion guarantee -

"95% of traffic matrices will receive enough capacity"

CMP (Chip Multi-Processor)



How can we model the future traffic matrix distribution?

Traffic Matrix Distribution

- Measure on real large CMP networks?
- Problems
 - > Their future applications are unknown
 - Their future traffic types are unknown: Between processors? Cache accesses? Control traffic?
 - Chicken-and-egg problem: traffic might depend on what architecture offers)

Traffic Matrix Distribution

- General model: any core can communicate with any core - but with a bounded core input/output rate.
 - > Example: each core may send/receive up to 1 Gbps.
 - Used in CMPs [Murali et al. '07] but also backbone networks [Dukkipati et al. '05], interconnection networks [Towles and Dally '02], VPN networks [Duffield et al. '99], routers [McKeown et al. '96]

Core 1 sends
$$\longrightarrow$$

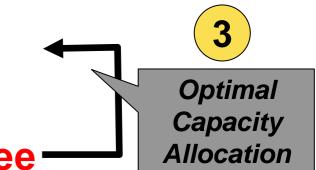
 $T = \begin{pmatrix} 0 & 0.5 & 0.1 & 0.2 \\ 0 & 0 & 0 & 0.6 \\ 0.2 & 0.1 & 0 & 0 \\ 0.1 & 0.1 & 0.3 & 0 \end{pmatrix} \leq 1$

Reminder: just one model among many...

Statistical Approach to NoC Design

Given:

- Traffic matrix distribution
 - Topology
 - Routing
 - Link capacities



Compute congestion guarantee

"95% of traffic matrices will receive enough capacity"

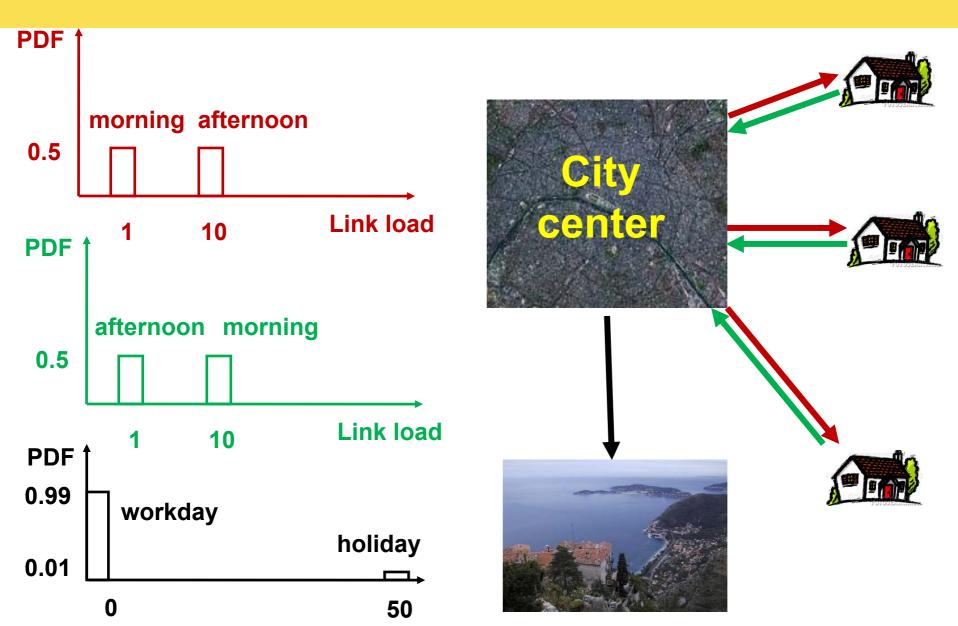
Compute Congestion Guarantee

Show that "95% of all traffic matrices will receive enough capacity on link *l*"

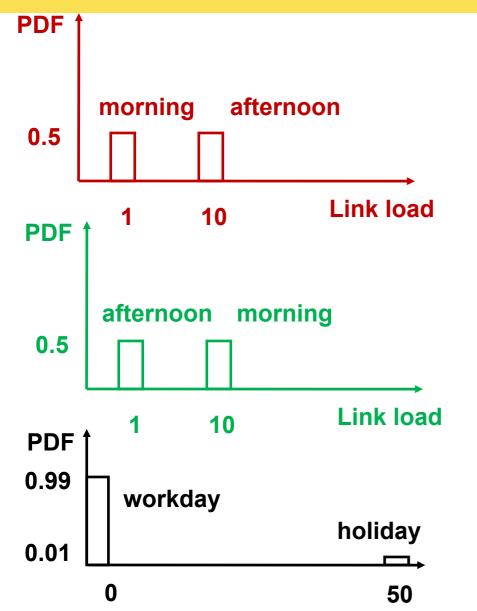
\Leftrightarrow

- 1. Compute *Traffic-load distribution Plot* (or *T-Plot*) for link *(*
- 2. Show that the load on link *ℓ* is less than its capacity for **95%** of traffic matrices.

Link T-Plot

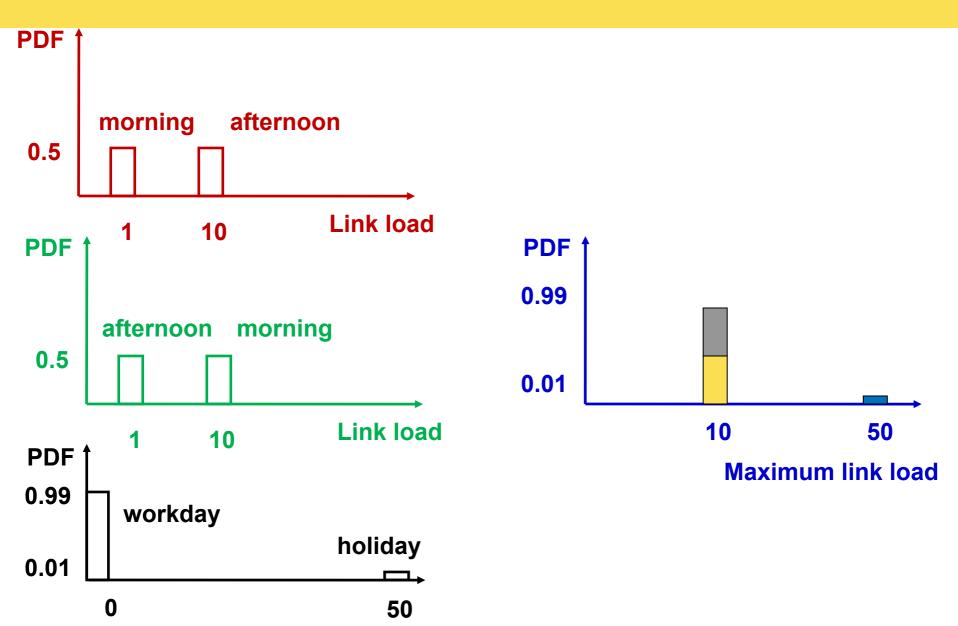


Global T-Plot

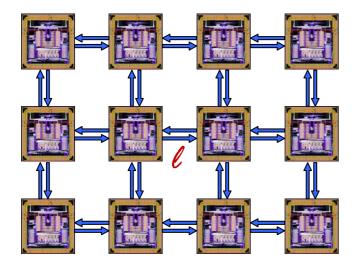


- We want to provide statistical guarantees on the whole network
 - i.e. on all links
- Global T-Plot: for each traffic matrix *T*, measure maximum load among all links





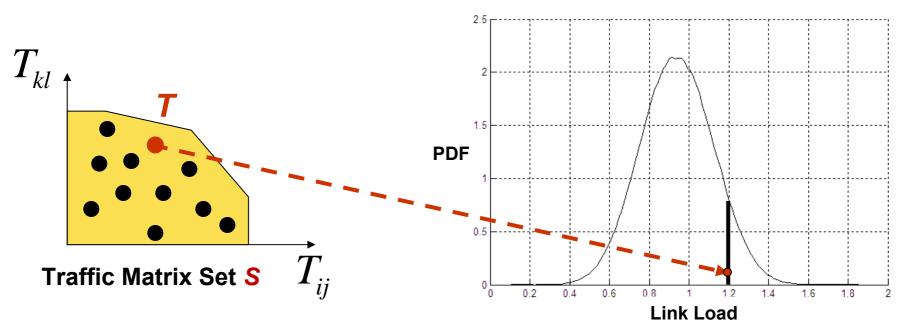
Local T-Plots in NoCs



Given:

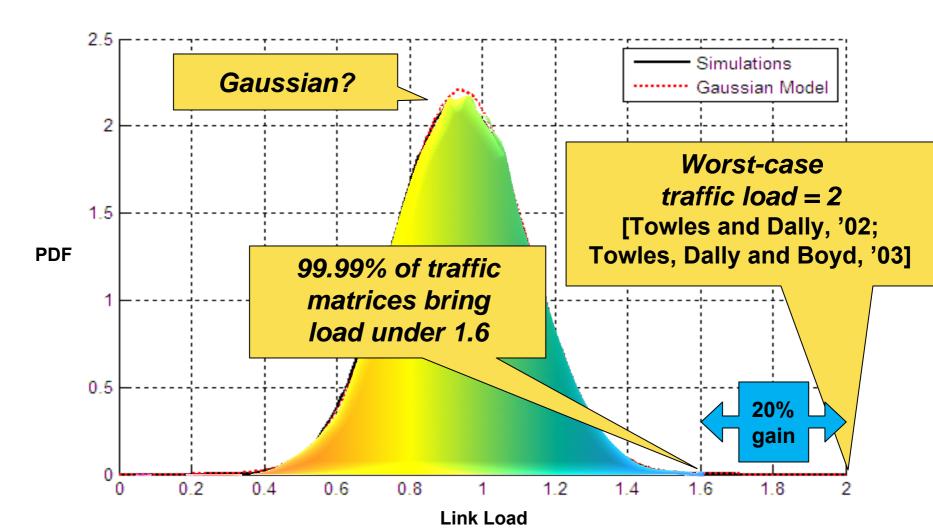
- ➤ Traffic matrices T∈S
- > 3x4 mesh topology
- XY routing
- Link l

Find T-Plot on



T-Plot

Close-up view:

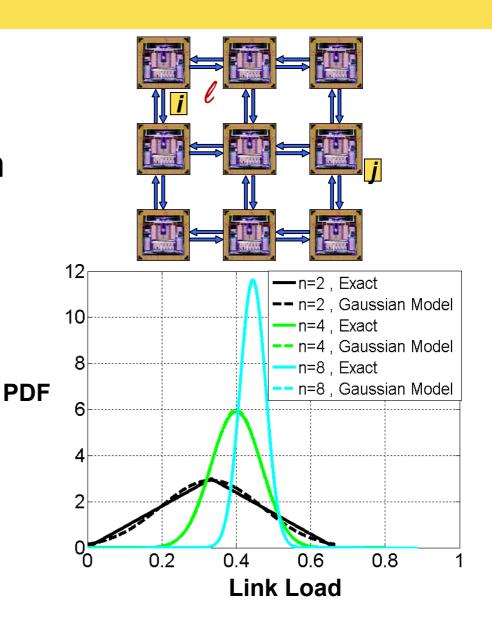


Are T-Plots Gaussian?

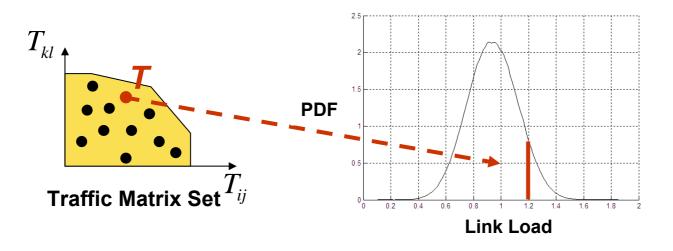
- Some T-Plots may be well approximated as Gaussian.
- Intuition (Central Limit Theorem): when N grows, the sum of N i.i.d. (independent and identically distributed) random variables converges to a Gaussian distribution.

Gaussian T-Plot Example

- Example: n x n mesh with XY routing.
- Assume i.i.d traffic from *i* to *j*
- Theorem: As n grows, T-Plot for converges to Gaussian.

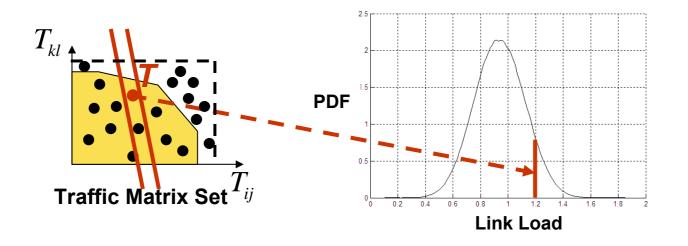


Computing the T-Plot



- Theorem: for an arbitrary graph and routing, computing the *T-Plot* is *#P-complete*.
- #P-complete problems are at least as hard as NP-complete problems.
 - > NP: "Is there a solution?"
 - #P: "How many solutions?"

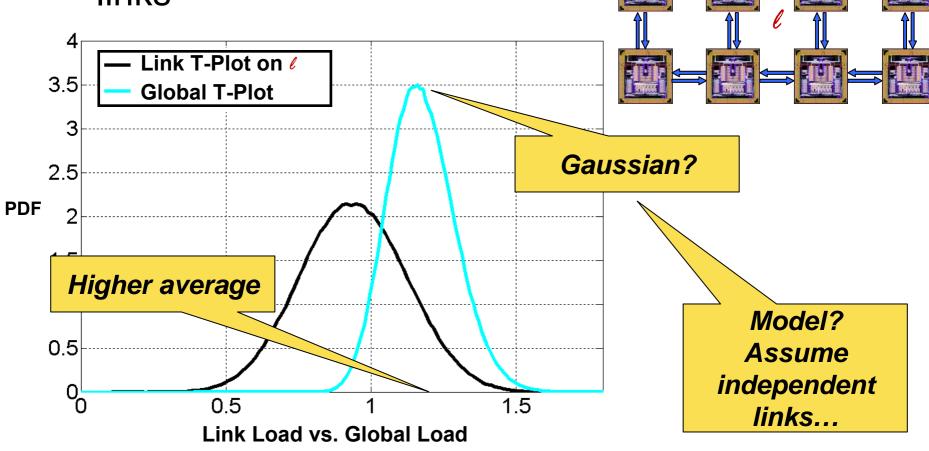
Computing the T-Plot



- Idea: use Monte Carlo simulations to approximate T-Plots
 - Plot many points, and count number of points to approximate density.

Link T-Plot Vs Global T-Plot

Global T-Plot: for each traffic matrix T, measure maximum load among all links



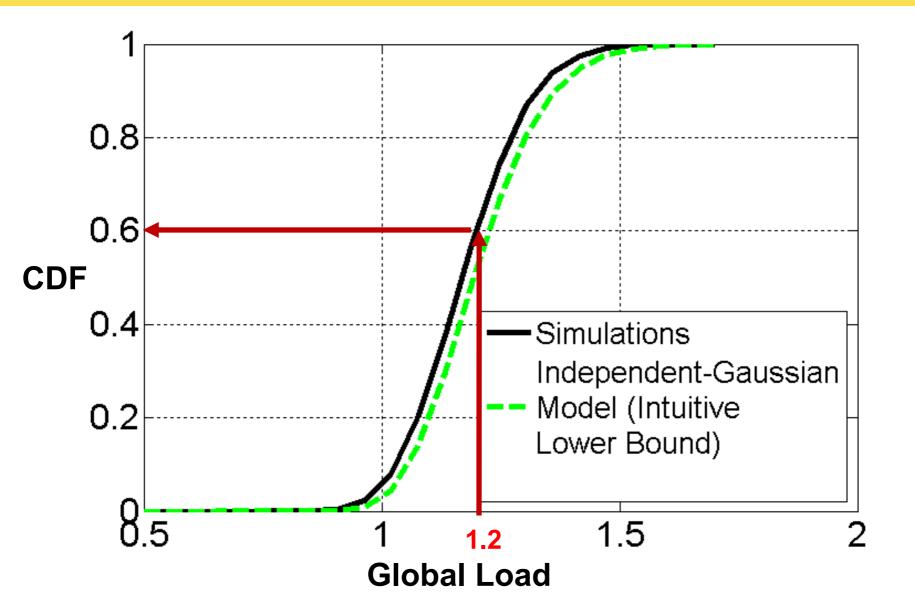
Link-Independent Model: Simple Example

> Assume:

> If ℓ_1 and ℓ_2 independent:

Global Load:
$$\begin{cases} 0 & w.p.\frac{1}{4} \\ 1 & w.p.\frac{3}{4} \end{cases}$$

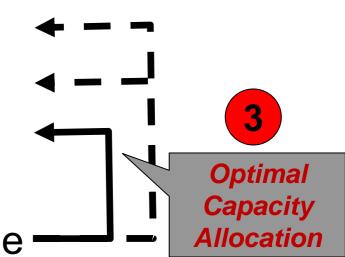
Independent-Gaussian Model



Statistical Approach to NoC Design

Given:

- Traffic matrix distribution
 - Topology
 - Routing
 - Link capacities

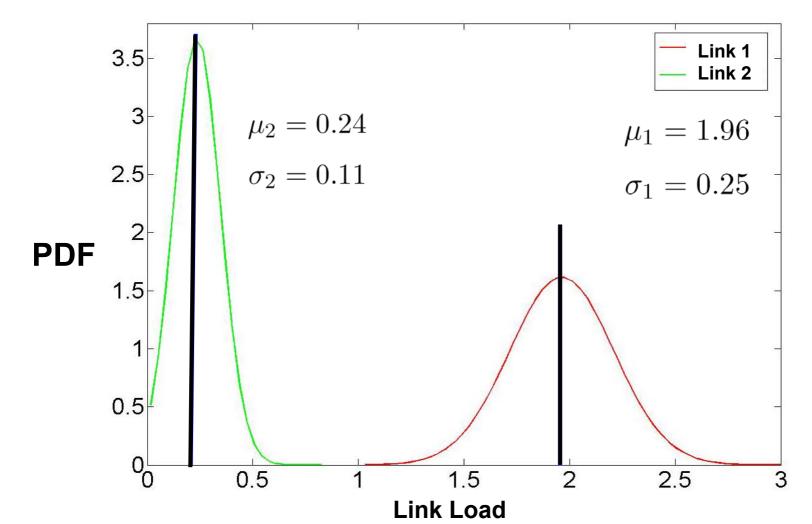


Compute congestion guarantee –

"95% of traffic matrices will receive enough capacity"

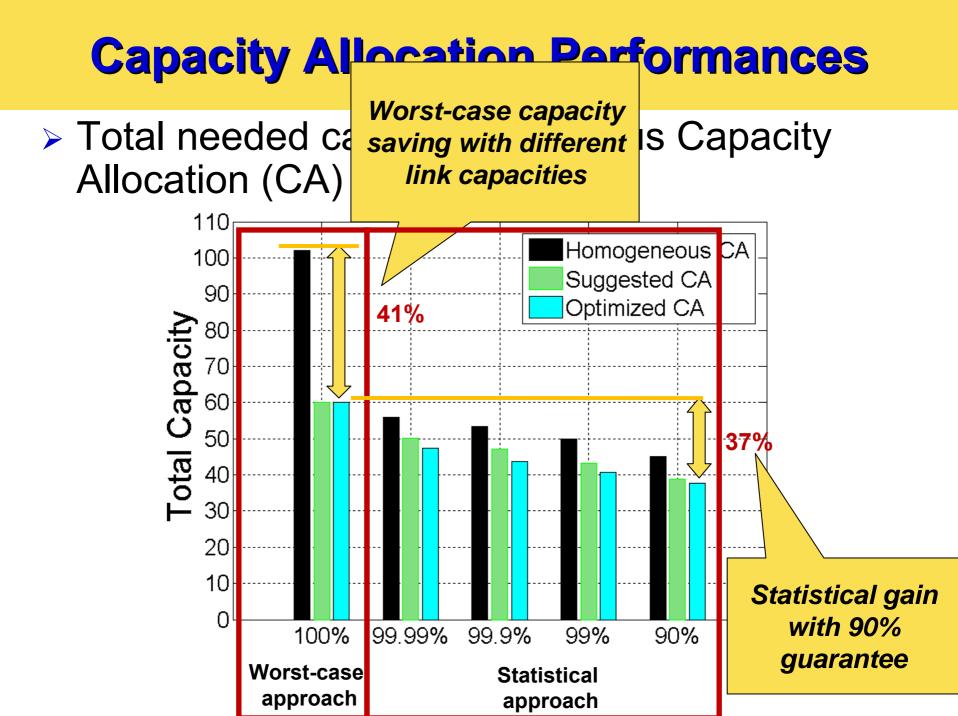
Capacity Allocation Intuition

Idea: given a total capacity, distribute it so that the overflow probability on each link is the same.



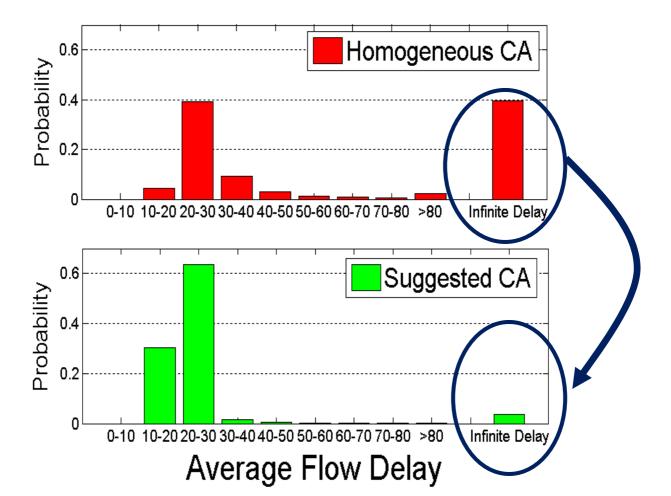
Capacity Allocation Intuition

- Idea: given a total capacity, distribute it so that the overflow probability on each link is the same.
- Theorem: if Gaussian-independent model holds, with same standard-deviation σ on all links, then this scheme is optimal.

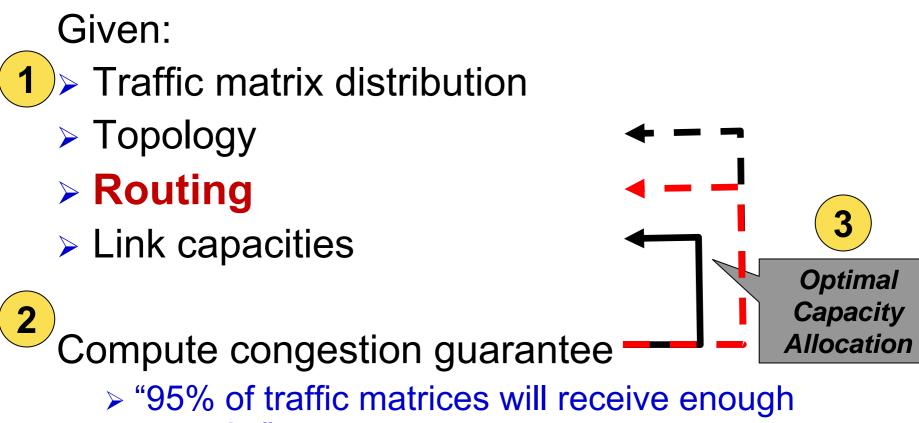


Average Flow Delay Model

- > Assume M/M/1 delay model
- Average flow delay distribution over all traffic matrices, for two different Capacity Allocation (CA) schemes:



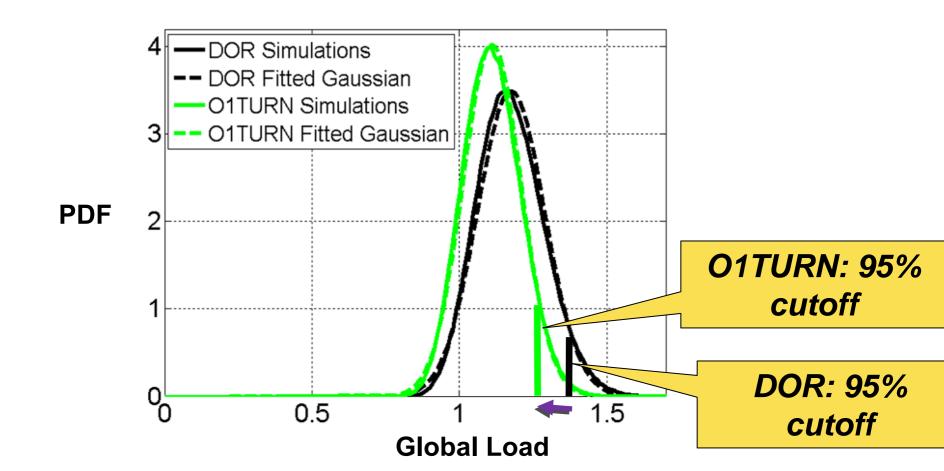
Statistical Approach to NoC Design



capacity"

Comparing routing algorithms

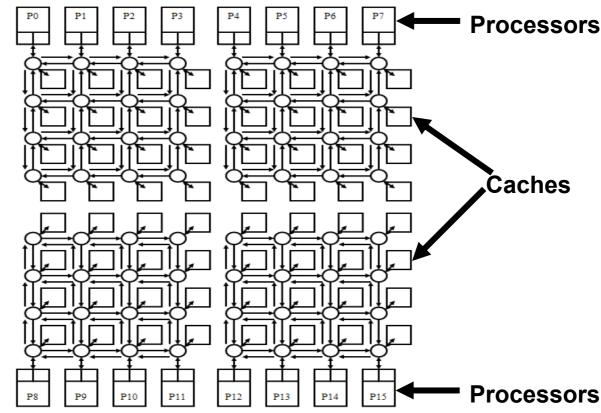
> DOR (XY) and O1TURN ($\frac{1}{2}$ XY, $\frac{1}{2}$ YX) on 3x4 mesh:



NUCA network

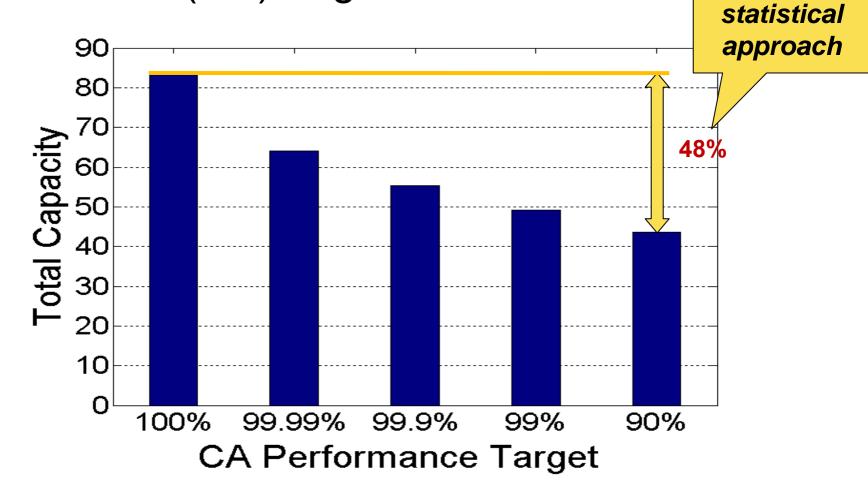
More complex CMP architectures show similar results

- NUCA (Non-Uniform Cache Architecture) with sharing degree 4.
- Traffic model: each core (cache) may only send/receive traffic to/from caches (cores) in its sub-network.



NUCA network – Total capacity

Total capacity required for various Capacity Allocation (CA) targets.



Summary

- Statistical approach to NoC design
 - > Deals with several traffic matrices
 - Can optimize capacity allocation and routing
- Can be applied to any
 - > Traffic matrix distribution
 - > Topology
 - > Oblivious routing
 - Link capacities

Thank you.